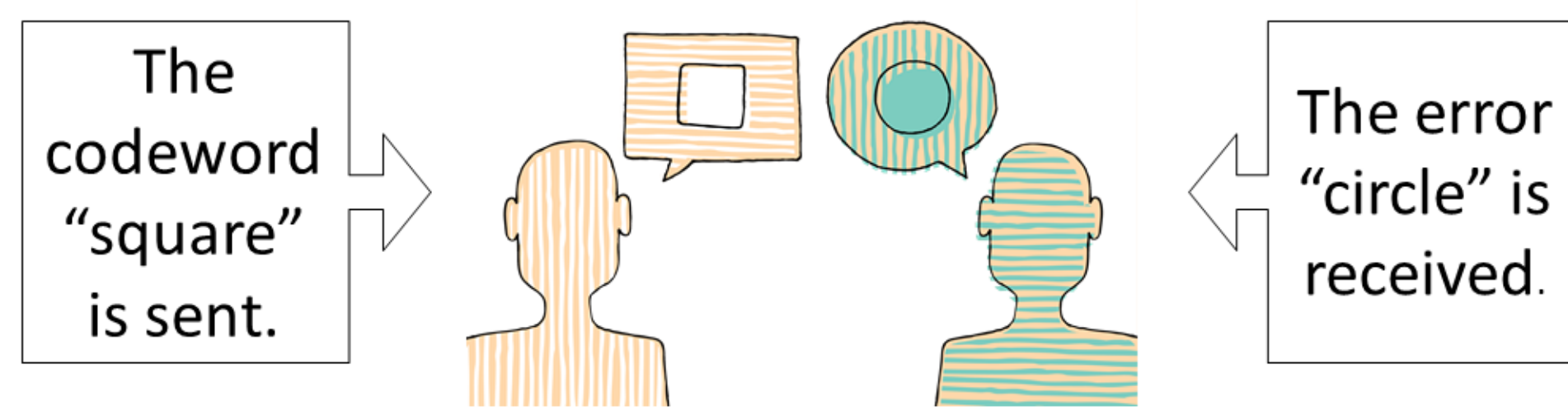


# Error Correction in Finite Arbitrary Length Binary Messages

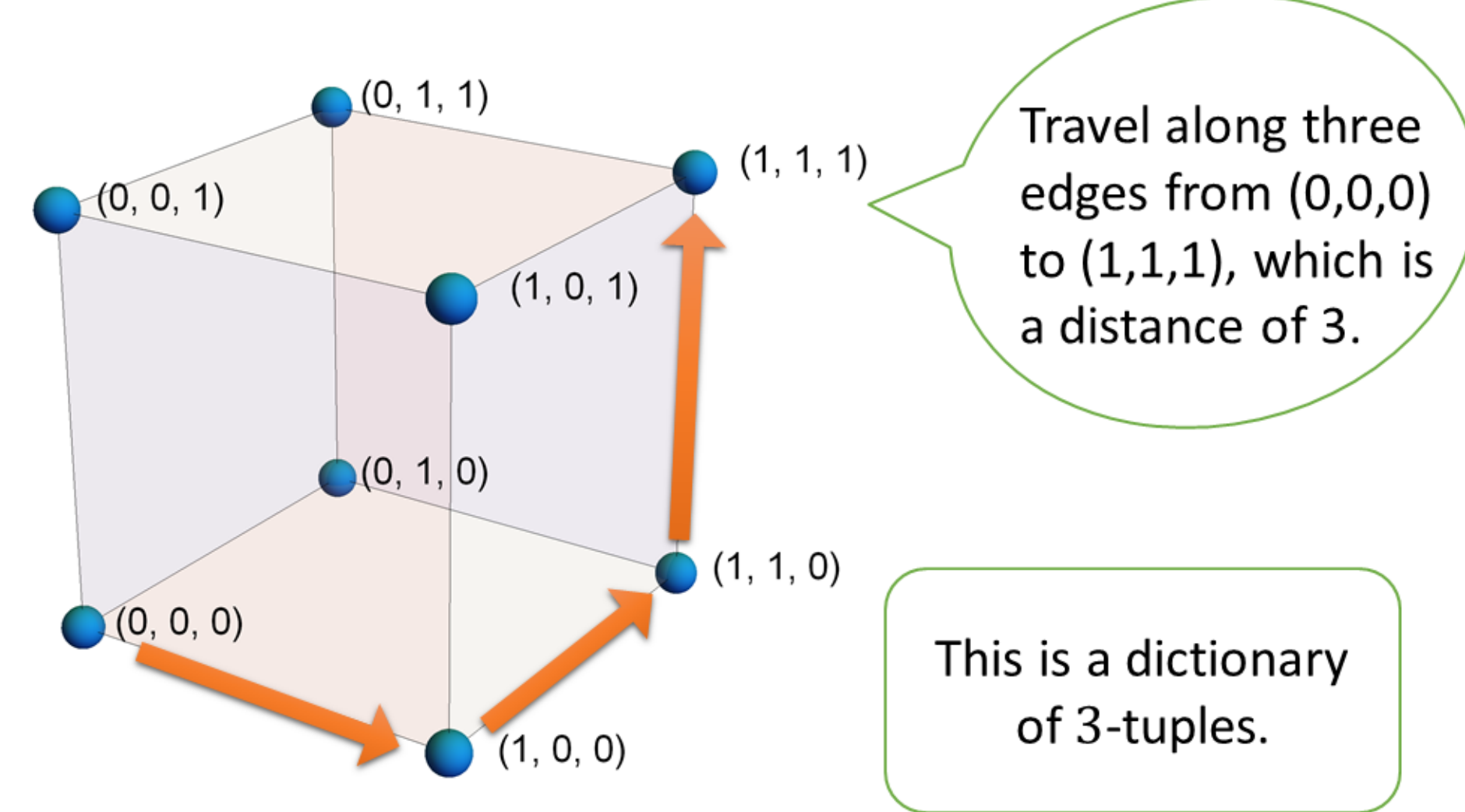
Dr. Donald White, Wayne Fincher, Jeremy Bouza, and special thanks to Dr. Aron

## Efficient Digital Communications



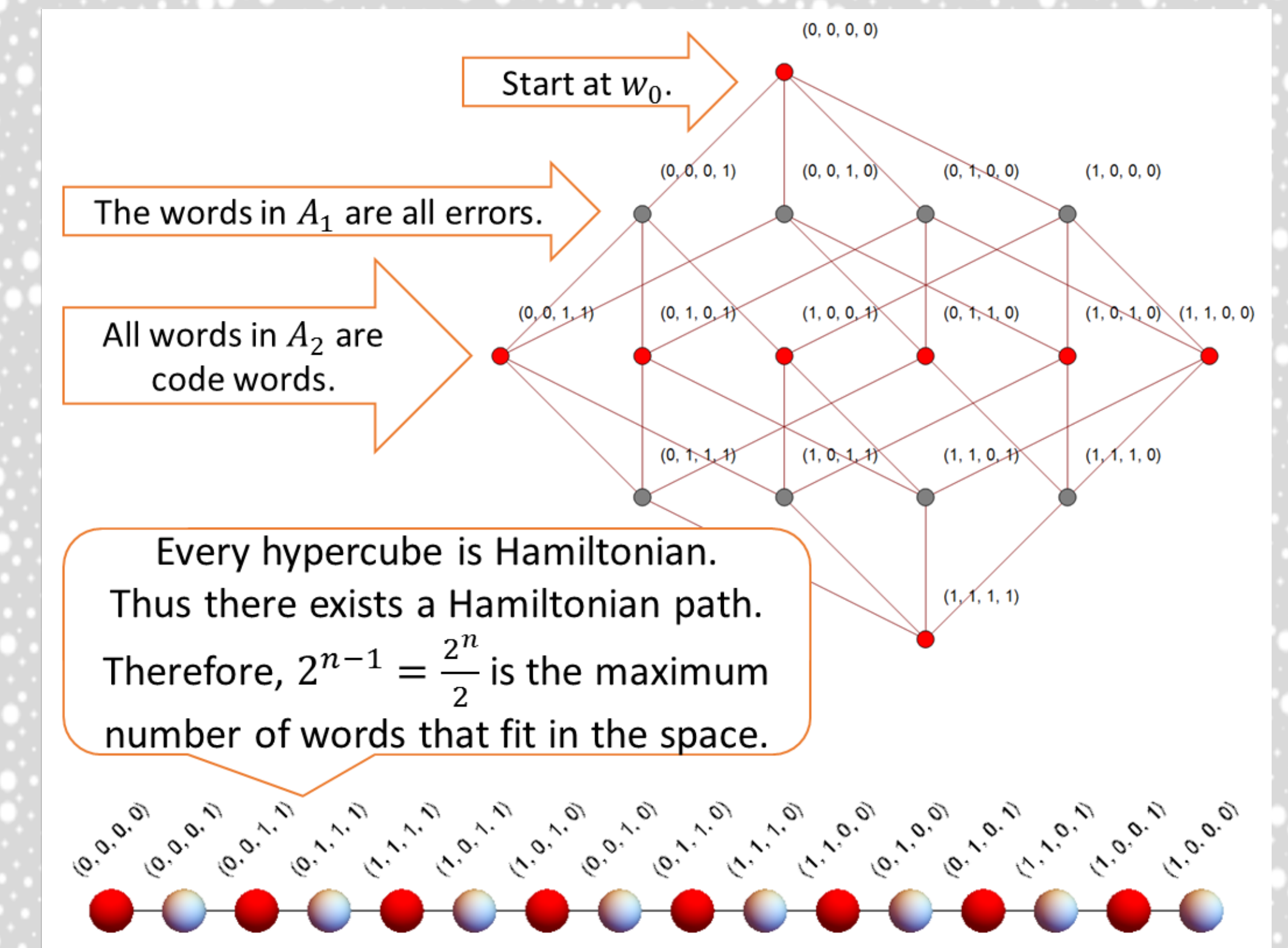
- 1<sup>ST</sup> Can errors be detected?
- 2<sup>ND</sup> Can errors be corrected?

## The 3-Dimensional Picture



## How Many Words Are In a Ball?

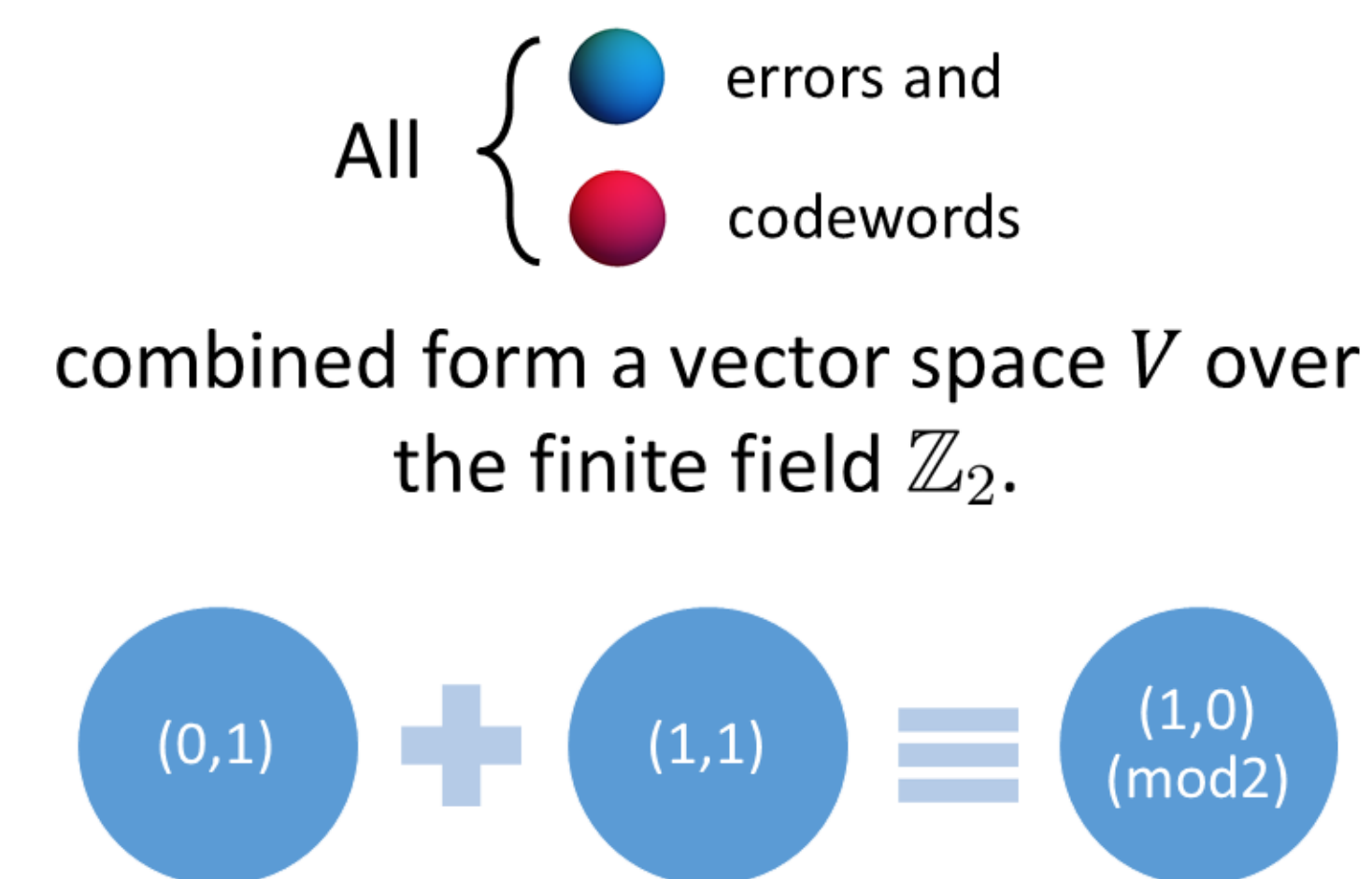
- The image shows all of the words in a dictionary consisting of 4-tuples.
- From one level to the next are words with a distance of one.
- The pattern is that of binomial coefficients.
- The general formula for the number of words in the  $B_H(w, j)$  where  $1 \leq j \leq n$  is
 
$$\sum_{k=0}^j \binom{n}{k}.$$



## What are the codewords?

- Dictionary: A vector space  $V$  over the field  $\mathbb{Z}_2$ .
- Word: A vector in  $V$ . (i.e., a codeword or an error)
- Distance: The number of components that are different between two words.

## The Abstract Picture



## Words in an Annulus Have Even Distances

**Theorem:** Let  $V$  be a vector space over  $\mathbb{Z}_2$ . For any  $w_0 \in V$ , the set  $A_j = \{y: H(w_0, y) = j, \forall y \in V\}$  does not contain any two words an odd distance apart.

*Proof.* Starting from  $w_0$  words that are a distance of one from  $w_0$ , denoted  $A_1$ , each have one component that is different from  $w_0$ . Each of those words in  $A_1$  has two components different from the other words in  $A_1$ .

Assume that the words in  $A_j$  all have an even distance from the other words in  $A_j$ . Let us consider the set  $A_{j+1}$ , and let  $a$  be an element of  $A_{j+1}$  and let  $a'$  be the element in  $A_j$  such that one component of  $a'$  is transposed to obtain  $a$ . By symmetry, we only need to consider the distance between  $a$  and the elements of  $\{A_{j+1} - a\}$ , which we shall denote as  $B$ .

Consider the distance between  $a$  and some element  $b \in B$  such that  $b' \in A_j$  and one component of  $b'$  is transposed to obtain  $b$ . If the component of  $a'$  that is transposed is the same component of  $b'$  that is transposed, then  $H(a, b) = H(a', b')$ . Otherwise,  $H(a, b) = H(a', b') + 2$ .  $\square$

## Maximum Correctable Words

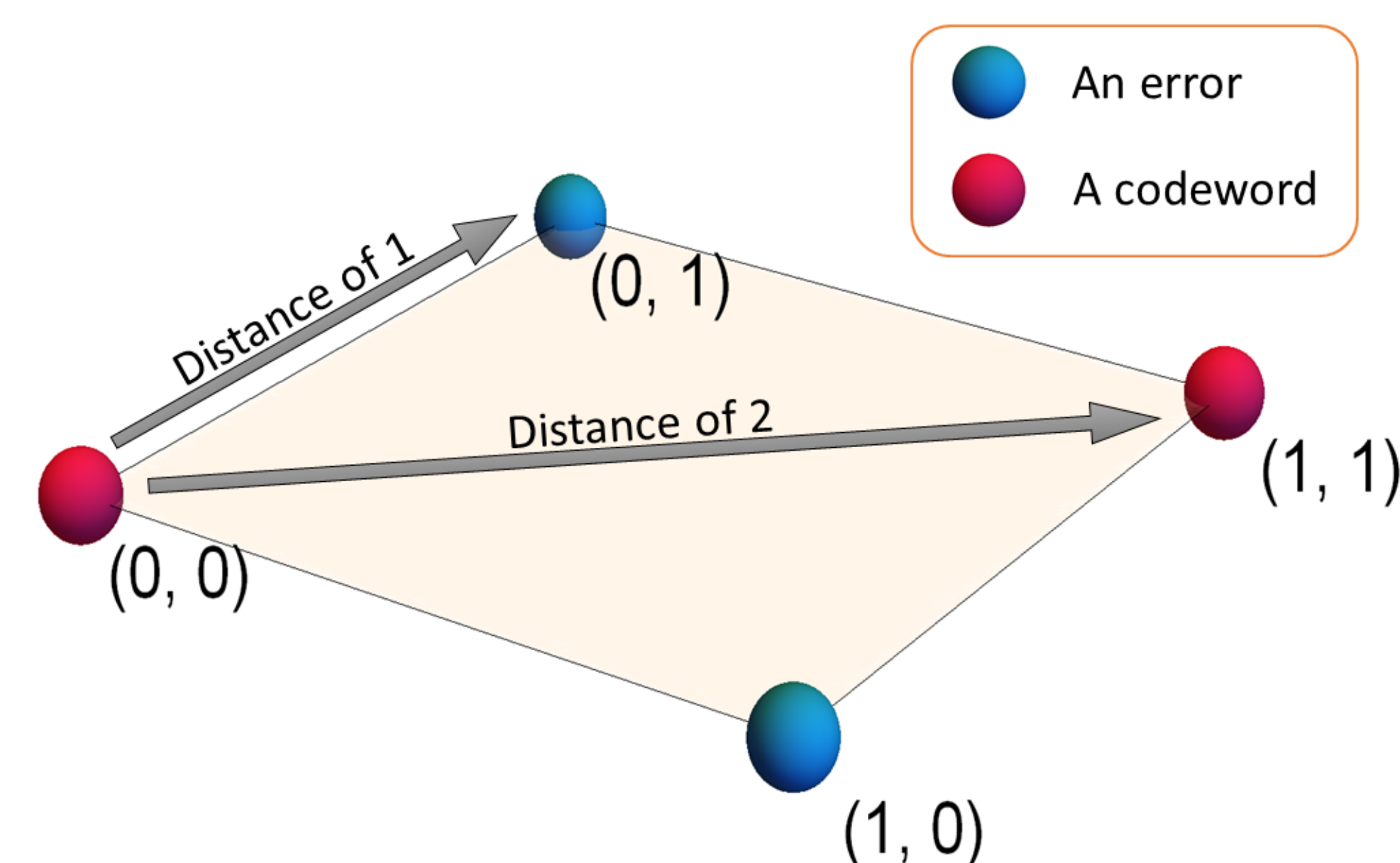
**Lower bound:** Assuming that the ball of radius 2 centered at any codeword is disjoint for any other similar ball, then the lower bound is:

$$\left\lfloor \frac{2^{n+1}}{n^2 + n + 1} \right\rfloor$$

**Upper bound:** Since there is a path connecting all words in the dictionary in a single line, at most every third word can be a codeword:

$$\frac{2^n}{3}$$

## A Geometric Interpretation



## Symmetry of the Space

Note: We can always let (0,0,...,0) be the first codeword because the space is symmetric.

*Proof.* Let  $c'$  be some codeword in  $V$  and  $f: V \rightarrow V$  be defined as

$$f(c) = \left( \left\{ \begin{array}{l} c_1, c'_1 = 0 \\ c_1 + 1, c'_1 = 1 \end{array} \right\}, \dots, \left\{ \begin{array}{l} c_n, c'_n = 0 \\ c_n + 1, c'_n = 1 \end{array} \right\} \right).$$

It is clear that  $f(f(c)) = c$  for any  $c \in V$ . Thus,  $f$  is its own inverse function. A map from one discrete space to another discrete space always maps open sets to open sets. Therefore  $f$  is a homeomorphism.  $\square$

## Maximum Detectable Code Words

- Pick a  $w_0$  to be the first code word.
- No words in the annulus  $A_1$  can be a code word.
- We showed that all words in  $A_2$  are at least a distance of two from each other. Chose all of them as code words.
- Every word in  $A_3$  is a distance of one from some code word in  $A_2$ , so there are no code words here.
- Keep choosing every word in every other annulus.
- This gives  $2^{n-1}$  detectable code words.
- But can we fit even more code words into the space?

## Conclusions

- The number of single-error correctable codewords in the space  $n = 4$  is 2. For  $n = 5$ , it is 4, and for  $n = 6$ , we have 8. Therefore, our conjecture is that there are  $2^{n-3}$  codewords in the space.
- Our future research will focus on narrowing the upper and lower bounds of the maximum correctable code words until they converge.
- Additionally, we are investigating the use of generator and parity check matrices in order to

## Acknowledgements

- We thank Dr. Aron and Dr. White for their guidance.
- Colin Adams, and Robert Franzosa. *Introduction to Topology: Pure and Applied*. Upper Saddle River, NJ: Pearson Prentice Hall, 2008. 150-152.
- Sarah Adams, *Introduction to Algebraic Coding Theory*. <http://www.math.cornell.edu/~web3360/eccbook2007.pdf>, 2008.